



Linux Ubuntu 26.04 Manual Pages on command 'pthread_cond_destroy.3'

\$ man pthread_cond_destroy.3

pthread_cond_init(3) Library Functions Manual pthread_cond_init(3)

NAME

pthread_cond_init, pthread_cond_signal, pthread_cond_broadcast, pthread_cond_wait,
pthread_cond_timedwait, pthread_cond_destroy - operations on conditions

SYNOPSIS

```
#include <pthread.h>

pthread_cond_t cond = PTHREAD_COND_INITIALIZER;

int pthread_cond_init(pthread_cond_t *cond,
                     pthread_condattr_t *cond_attr);

int pthread_cond_signal(pthread_cond_t *cond);

int pthread_cond_broadcast(pthread_cond_t *cond);

int pthread_cond_wait(pthread_cond_t *cond, pthread_mutex_t *mutex);

int pthread_cond_timedwait(pthread_cond_t *cond, pthread_mutex_t *mutex,
                          const struct timespec *abstime);

int pthread_cond_destroy(pthread_cond_t *cond);
```

DESCRIPTION

A condition (short for "condition variable") is a synchronization device that allows threads to suspend execution and relinquish the processors until some predicate on shared data is satisfied. The basic operations on conditions are: signal the condition (when the predicate becomes true), and wait for the condition, suspending the thread execution until another thread signals the condition.

A condition variable must always be associated with a mutex, to avoid the race condition where a thread prepares to wait on a condition variable and another thread signals the con?

dition just before the first thread actually waits on it.

`pthread_cond_init()` initializes the condition variable `cond`, using the condition attributes specified in `cond_attr`, or default attributes if `cond_attr` is `NULL`. The LinuxThreads implementation supports no attributes for conditions, hence the `cond_attr` parameter is actually ignored.

Variables of type `pthread_cond_t` can also be initialized statically, using the constant `PTHREAD_COND_INITIALIZER`.

`pthread_cond_signal()` restarts one of the threads that are waiting on the condition variable `cond`. If no threads are waiting on `cond`, nothing happens. If several threads are waiting on `cond`, exactly one is restarted, but it is not specified which.

`pthread_cond_broadcast()` restarts all the threads that are waiting on the condition variable `cond`. Nothing happens if no threads are waiting on `cond`.

`pthread_cond_wait()` atomically unlocks the mutex (as per `pthread_unlock_mutex()`) and waits for the condition variable `cond` to be signaled. The thread execution is suspended and does not consume any CPU time until the condition variable is signaled. The mutex must be locked by the calling thread on entrance to `pthread_cond_wait()`. Before returning to the calling thread, `pthread_cond_wait()` re-acquires mutex (as per `pthread_mutex_lock()`).

Unlocking the mutex and suspending on the condition variable is done atomically. Thus, if all threads always acquire the mutex before signaling the condition, this guarantees that the condition cannot be signaled (and thus ignored) between the time a thread locks the mutex and the time it waits on the condition variable. See CAVEATS below.

`pthread_cond_timedwait()` atomically unlocks mutex and waits on `cond`, as `pthread_cond_wait()` does, but it also bounds the duration of the wait. If `cond` has not been signaled within the amount of time specified by `abstime`, the mutex is re-acquired and `pthread_cond_timedwait()` returns the error `ETIMEDOUT`. The `abstime` parameter specifies an absolute time, with the same origin as `time(2)` and `gettimeofday(2)`: an `abstime` of 0 corresponds to 00:00:00 GMT, January 1, 1970.

`pthread_cond_destroy()` destroys a condition variable, freeing the resources it might hold. No threads must be waiting on the condition variable on entrance to `pthread_cond_destroy()`. In the LinuxThreads implementation, no resources are associated with condition variables, thus `pthread_cond_destroy()` actually does nothing except checking that the condition has no waiting threads.

`pthread_cond_wait()` and `pthread_cond_timedwait()` are cancellation points. If a thread is cancelled while suspended in one of these functions, the thread immediately resumes execution, then locks again the mutex argument to `pthread_cond_wait()` and `pthread_cond_timedwait()`, and finally executes the cancellation. Consequently, cleanup handlers are assured that mutex is locked when they are called.

ASYNCR-SIGNAL SAFETY

The condition functions are not asynchr-signal safe, and should not be called from a signal handler. In particular, calling `pthread_cond_signal()` or `pthread_cond_broadcast()` from a signal handler may deadlock the calling thread.

RETURN VALUE

All condition variable functions return 0 on success and a non-zero error code on error.

ERRORS

`pthread_cond_init()`, `pthread_cond_signal()`, `pthread_cond_broadcast()`, and `pthread_cond_wait()` never return an error code.

The `pthread_cond_timedwait()` function returns the following error codes on error:

ETIMEDOUT

The condition variable was not signaled until the timeout specified by `abstime`.

The `pthread_cond_destroy()` function returns the following error code on error:

`EBUSY` Some threads are currently waiting on cond.

SEE ALSO

`pthread_condattr_init(3)`, `pthread_mutex_lock(3)`, `pthread_mutex_unlock(3)`, `gettimeofday(2)`, `nanosleep(2)`.

CAVEATS

The implementation of the provided functions until glibc 2.25 used an internal data lock.

This lock did not support priority-inheritance and was subject to unbounded priority inversion, visible on a real-time system.

After the rewrite of the implementation in glibc 2.25 the usage of internal lock changed.

The internal lock is always acquired by the signaling functions `pthread_cond_signal()` and `pthread_cond_broadcast()`. The waiting function acquires the lock if the waiting process was interrupted. The interruption can be caused for instance by a specified timeout, and denoted by the error value `ETIMEDOUT`, or by a received signal, which is denoted by the error value `EINTR`.

EXAMPLE

Consider two shared variables `x` and `y`, protected by the mutex `mut`, and a condition variable `cond` that is to be signaled whenever `x` becomes greater than `y`.

```
int x,y;

pthread_mutex_t mut = PTHREAD_MUTEX_INITIALIZER;

pthread_cond_t cond = PTHREAD_COND_INITIALIZER;
```

Waiting until `x` is greater than `y` is performed as follows:

```
pthread_mutex_lock(&mut);

while (x <= y) {

    pthread_cond_wait(&cond, &mut);

}

/* operate on x and y */

pthread_mutex_unlock(&mut);
```

Modifications on `x` and `y` that may cause `x` to become greater than `y` should signal the condition if needed:

```
pthread_mutex_lock(&mut);

/* modify x and y */

if (x > y) pthread_cond_broadcast(&cond);

pthread_mutex_unlock(&mut);
```

If it can be proved that at most one waiting thread needs to be waken up (for instance, if there are only two threads communicating through `x` and `y`), `pthread_cond_signal()` can be used as a slightly more efficient alternative to `pthread_cond_broadcast()`. In doubt, use `pthread_cond_broadcast()`.

To wait for `x` to become greater than `y` with a timeout of 5 seconds, do:

```
struct timeval now;

struct timespec timeout;

int retcode;

pthread_mutex_lock(&mut);

gettimeofday(&now, NULL);

timeout.tv_sec = now.tv_sec + 5;

timeout.tv_nsec = now.tv_usec * 1000;

retcode = 0;

while (x <= y && retcode != ETIMEDOUT) {
```

```
        retcode = pthread_cond_timedwait(&cond, &mut, &timeout);
    }
    if (retcode == ETIMEDOUT) {
        /* timeout occurred */
    } else {
        /* operate on x and y */
    }
    pthread_mutex_unlock(&mut);
```